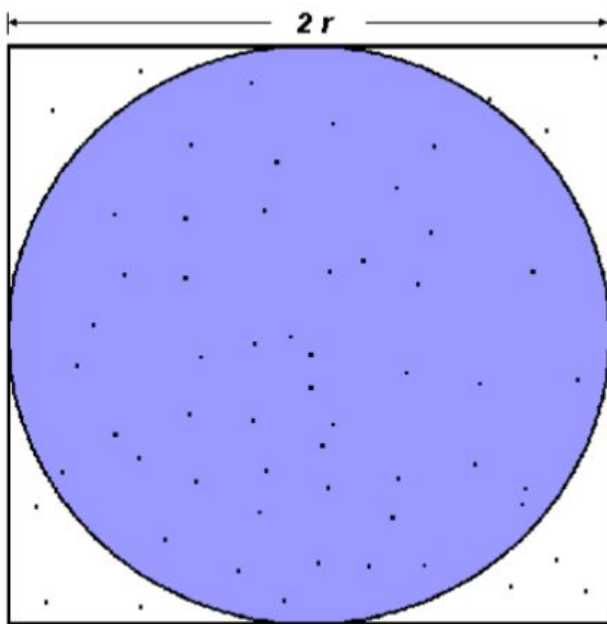


# URL sources

<https://github.com/lstorchi/mpisttest>

# Compute PI using a Monte Carlo Approach

A circle of radius  $R$  is inscribed inside a square with side length  $2 \cdot R$ , if so the area of the circle will be  $A_c = \pi R^2$  and the area of the square will be  $A_s = (2 \cdot R)^2$ . So the ratio of the area of the circle to the area of the square will be  $A_c / A_s = \pi/4$ .



$$A_s = (2r)^2 = 4r^2$$

$$A_c = \pi r^2$$

$$\pi = 4 \times \frac{A_c}{A_s}$$

If a program picks  $N$  points  $(x, y)$  at random inside the square. If a point is inside the circle (i.e. if  $x^2 + y^2 < R^2$ )  $M$  is incremented by one.

Thus finally:  $\pi = 4 \cdot M / N$

# Pseudo code

```
N = 2000
circle_count = 0
for i = 1 to N
    x = random value (0.0, 1.0)
    y = random value (0.0, 1.0)
    if x2 + y2 < 1.0
        circle_count = circle_count + 1
    endif
endfor
pi = 4 * (circle_count / N)
```

# Parallel version

- Each one of the  $P$  MPI processes will generate  $N/P$  random points (clue each process should use a different seed)
- You need to sum the final value of `circle_count` (you may need to use the `MPI_Reduce`)
- As in the serial code you may now estimate the PI value
- $N$  could be a command line argument

# Exercise

- Implement the serial version of the code starting from the pseudo code ( suggestion the number of random points can be a command line argument )
- Implement the parallel version and, depending on the number of cores of the VM you are using, calculate the speedup (**MPI\_Wtime()** Returns time in seconds since an arbitrary time in the past. **clock\_t clock(void)** returns the number of clock ticks elapsed since the program was launched. To get the number of seconds used by the CPU, you need to divide by **CLOCKS\_PER\_SEC**.)

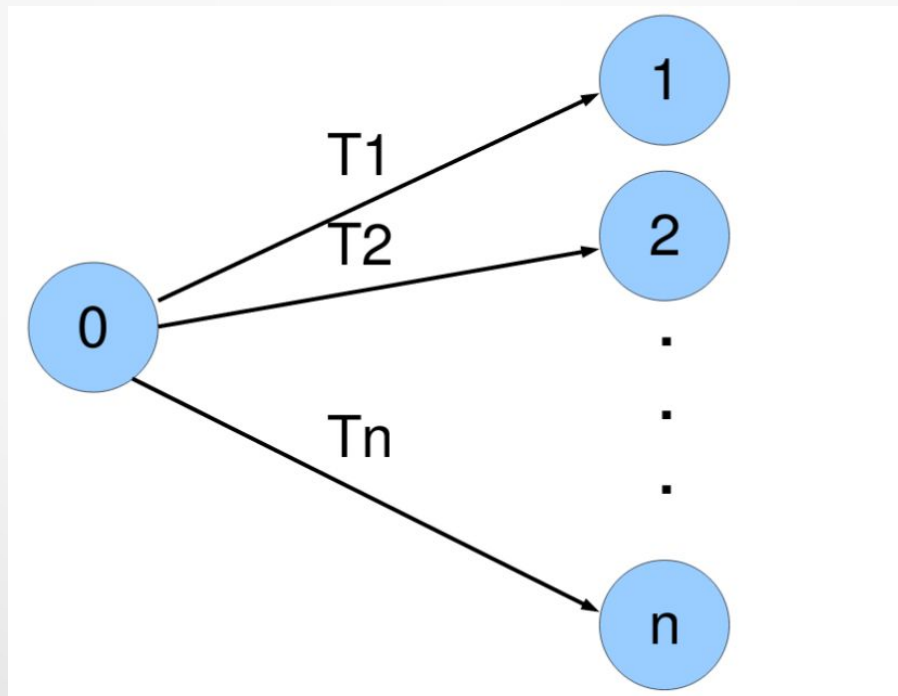
# MPI functions (C API)

- `MPI_Init (&argc, &argv);`
- `MPI_Comm_size (MPI_COMM_WORLD, &size);`
- `MPI_Comm_rank (MPI_COMM_WORLD, &rank);`
- `MPI_Barrier (MPI_COMM_WORLD);`
- `MPI_Reduce (&circle_count, &t_circle_count, 1, MPI_INT, MPI_SUM, 0, MPI_COMM_WORLD);`
- `MPI_Finalize ();`

# Broadcast

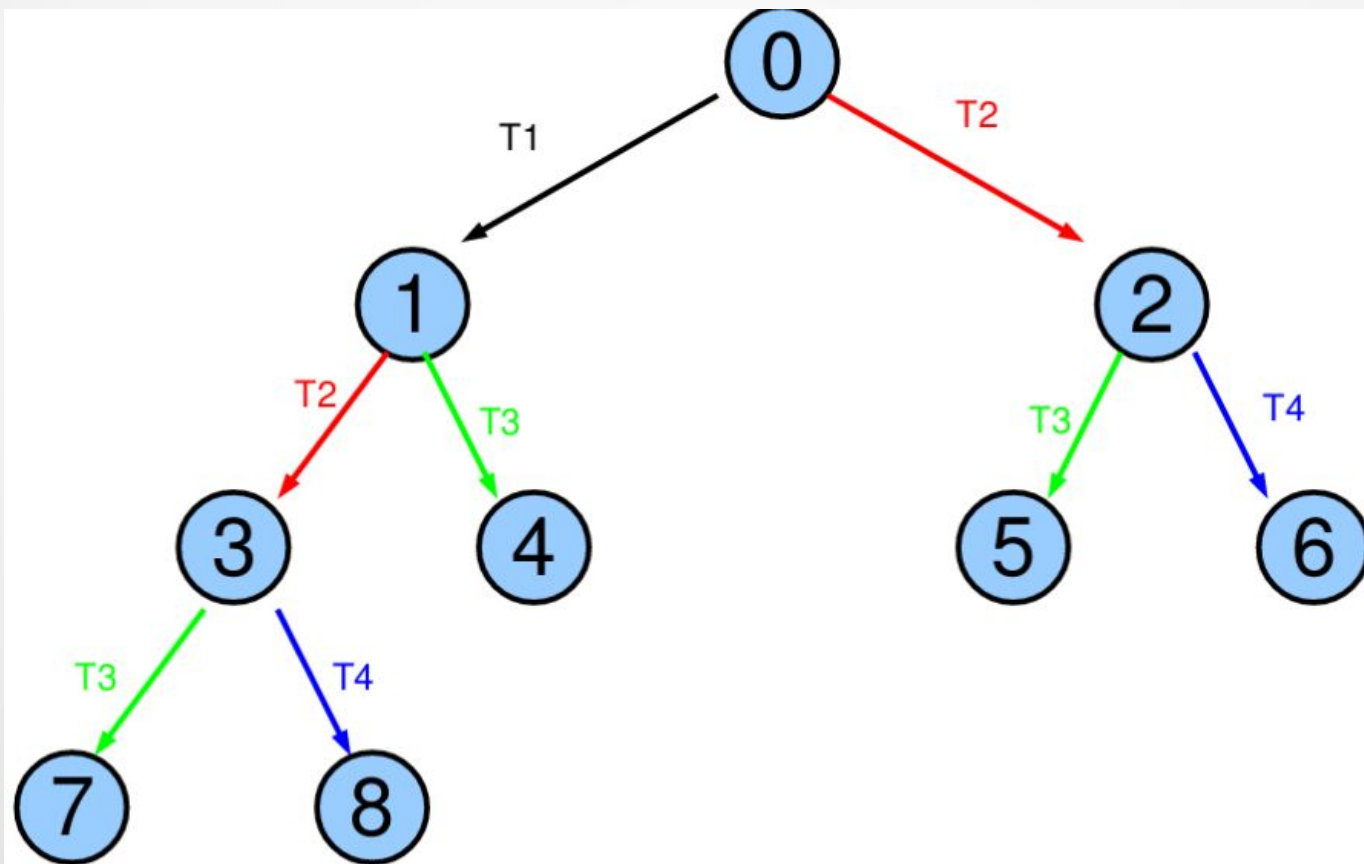
We will explore the differences between a naive approach to perform a broadcast and a more sophisticated one

Naive (flat tree) need  $n-1$  point to point communications,  $T(\text{msize})$  = time to send a message of size  $\text{msize}$ , so  $\text{Time} = (n-1) * T(\text{msize})$  so  $O(n)$



# Broadcast

Now using a binary tree we are able to reduce the communication time (Time =  $O(\log_2(n))$ )





# Bcast pseudocode

```
fromrank = (int) ((myrank-1)/2)
```

```
if (myrank > 0)
```

```
    Recv data from fromrank
```

```
torank1 = 2 * myrank + 1;
```

```
torank2 = 2 * myrank + 2;
```

```
if (torank1 < size)
```

```
    Send data to torank1
```

```
if (torank2 < size)
```

```
    Send data to torank2
```

# Exercise

- Implement the two version of the broadcast the one using the flat tree and the other using the binary tree. You will broadcast in both cases a vector of dimension  $N$ , where  $N$  again could be a command line argument.

# MPI functions (C API)

- `int MPI_Send(const void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm)`

```
MPI_Send (sbuf, bufdim, MPI_DOUBLE,
torank1, torank1, MPI_COMM_WORLD);
```

# MPI functions (C API)

- `int MPI_Recv(void *buf, int count, MPI_Datatype datatype, int source, int tag, MPI_Comm comm, MPI_Status *status)`

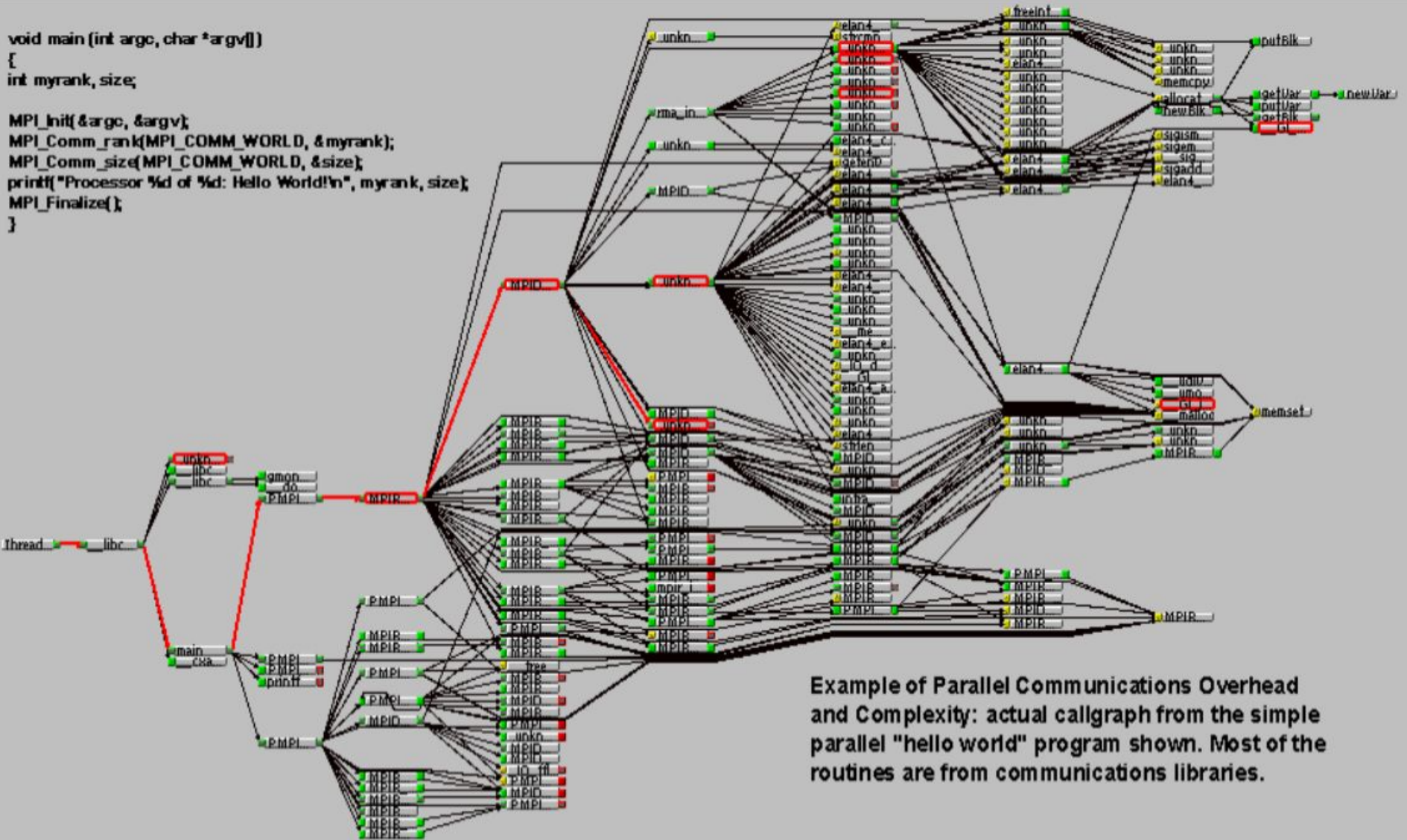
```
MPI_Status status;
```

```
MPI_Recv (sbuf, bufdim, MPI_DOUBLE,  
fromrank, myrank, MPI_COMM_WORLD,  
&status);
```

# MPI complexity

```
void main (int argc, char *argv[])
{
    int myrank, size;

    MPI_Init(&argc, &argv);
    MPI_Comm_rank(MPI_COMM_WORLD, &myrank);
    MPI_Comm_size(MPI_COMM_WORLD, &size);
    printf("Processor %d of %d: Hello World!\n", myrank, size);
    MPI_Finalize();
}
```



**Example of Parallel Communications Overhead and Complexity:** actual callgraph from the simple parallel "hello world" program shown. Most of the routines are from communications libraries.

# Serial code optimization

Parallel computing era, however .... to be cache friendly:

```
for (i=0; i<N; i++)  
  for (j=0; j<N; j++)  
    for (k=0; k<N; k++)  
      c[i][j] = c[i][j] + a[i][k] * b[k][j];
```

```
for (i=0; i<N; i++)  
  for (k=0; k<N; k++)  
    for (j=0; j<N; j++)  
      c[i][j] = c[i][j] + a[i][k] * b[k][j];
```



# Serial code optimization

```
[redo@banquo serialopt (master)]$ ./mm.1
Time to initialize 0.050375 s.
Time 10.007438 s.
Total time 10.057895 s.
```

```
Mflops -----> 213.512236
Check -----> 268364458.846206
```

```
[redo@banquo serialopt (master)]$ ./mm.3
Time to initialize 0.027267 s.
Time 2.971154 s.
Total time 2.998506 s.
```

```
Mflops -----> 716.184543
Check -----> 268364458.846206
```

```
[redo@banquo serialopt (master)]$ diff mm.1.c mm.3.c
37,38c37,38
<     for (j=0; j<N; j++)
<         for (k=0; k<N; k++)
---
>     for (k=0; k<N; k++)
>         for (j=0; j<N; j++)
```

# Serial code optimization

Keep the pipeline full, loop unrolling:

```
for (i=0; i<N; i++) {  
    for (k=0; k<N; k++) {  
        for (j=0; j<N; j +=8) {  
            c[i][j] = c[i][j] + a[i][k] * b[k][j];  
            c[i][j+1] = c[i][j+1] + a[i][k] * b[k][j+1];  
            c[i][j+2] = c[i][j+2] + a[i][k] * b[k][j+2];  
            c[i][j+3] = c[i][j+3] + a[i][k] * b[k][j+3];  
            c[i][j+4] = c[i][j+4] + a[i][k] * b[k][j+4];  
            c[i][j+5] = c[i][j+5] + a[i][k] * b[k][j+5];  
            c[i][j+6] = c[i][j+6] + a[i][k] * b[k][j+6];  
            c[i][j+7] = c[i][j+7] + a[i][k] * b[k][j+7];  
        }  
    }  
}
```



# MPI LIBRARY

## POINT-TO-POINT COMM.

Blocking

Standard: MPI\_Send

Synchronous: MPI\_Ssend

Buffered: MPI\_Bsend

RECEIVE: MPI\_Recv

Non Blocking

Standard: MPI\_Isend

Synchronous: MPI\_Issend

Buffered: MPI\_Ibsend

RECEIVE: MPI\_Irecv

# MPI LIBRARY

## COLLECTIVE COMM.

All to One

MPI\_Reduce

MPI\_Gather

MPI\_Gatherv

One to All

MPI\_Bcast

MPI\_Scatter

MPI\_Scatterv

All to All

MPI\_Barrier

MPI\_All\_reduce

MPI\_All\_gather